Using the Troika Clique Partitioning Algorithm for Analyzing the cluster dynamics of market index correlation networks

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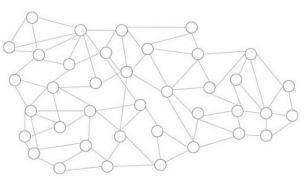
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Clique Partitioning (CP) Problem

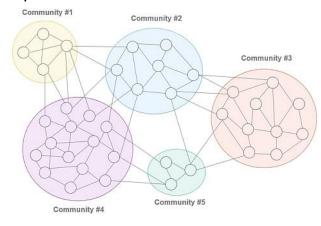
- <u>Input</u>: weighted undirected graph *G*(*V*,*E*) with *n* nodes and *m* edges
- <u>Method</u>: a clustering (clique partitioning) algorithm
- Output: a partition of the nodes into disjoint clusters (node colours)
- <u>Goal</u>: group nodes into clusters which maximizes the weight of the

partition



Output

Input



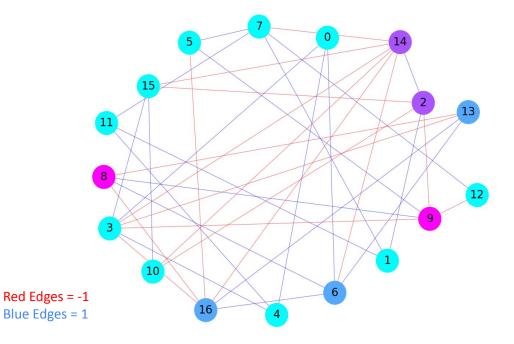
Figures from timbr.ai

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Grotschel & Wakabayashi (1989) A cutting plane algorithm for a clustering problem

How would you "clique partition" this graph?

<u>Goal</u>: Group nodes into clusters which maximizes the internal weight of the partition



<u>Maximized Objective</u> Objective value = 4 Value = 16

Investment Portfolio Analogy

Objective:

Maximize returns from a portfolio consisting of stocks and bonds.

Investments:

- Stocks
- Bonds

Variables:

- s = Amount invested in stocks
- b = Amount invested in bonds



Integer Programming (IP) Formulation

Maximize Returns:

Returns = 1.2s + 1.1b

Constraints:

1. Total investment limit:

 $s+b \leq 100,000$

2. Diversification:

 $s \ge 0, \quad b \ge 20000$

3. Integer constraint:

 $s, b \in \mathbb{Z}$



Optimal Solution

Solution:

Optimal investment in stocks (s): \$80,000

Optimal investment in bonds (b): \$20,000

Calculated Returns:

 $\mathsf{Max}\;\mathsf{Returns} = 1.2 \times 100000 + 1.1 \times 20000 = \$116,000$

Remarks:

- This allocation maximizes the returns under the given constraints.
- Preference is given to stocks due to higher returns per dollar invested.



Clique Partitioning IP Formulation

Given partition P, for every pair of nodes (i, j), their cluster assignment is either same (represented by xij = 0) or different (represented by xij = 1). The weight of the undirected edge between nodes (i, j) is represented as wij.

$$\operatorname{RP}^*(G): \max_{x_{ij}} W = \sum_{(i,j)\in E} w_{ij}(1-x_{ij})$$

Constraints:

s.t.
$$x_{ik} + x_{jk} \ge x_{ij} \quad \forall (i, j, k) \in T_{+}^{k}$$

 $x_{jk} + x_{ij} \ge x_{ik} \quad \forall (i, j, k) \in T_{+}^{j}$
 $x_{ij} + x_{ik} \ge x_{jk} \quad \forall (i, j, k) \in T_{+}^{i}$
 $x_{ij} \in \{0, 1\} \quad \forall (i, j) \in E$
 $T_{+}^{k} = \{(i, j, k) \in T \mid w_{ik} > 0 \lor w_{ij} > 0\}$
 $T_{+}^{i} = \{(i, j, k) \in T \mid w_{ij} > 0 \lor w_{ik} > 0\}.$



Existing Methods

Exact Methods:

- Solving IP formulation using commercial solvers (Miyauchi et al. 2018)
- Exact branch-and-cut exploring facet inequalities (Simanchev et al. 2019)

Heuristic Methods:

- Combo (Sobolevsky et al. 2019)
- Merge-divide memetic clique partitioning algorithm (Lu et al. 2021)



Troika

Inspiration: Bayan algorithm for Community Detection (Aref et al. 2022)





Troika

Branch-and-Cut Implementation

- Node Triple Method: Branch on sets of three nodes to explore feasible space.
- **Upper Bound**: Use linear programming relaxation to establish upper bounds.
- Lower Bound: Employ heuristic search to construct lower bounds.



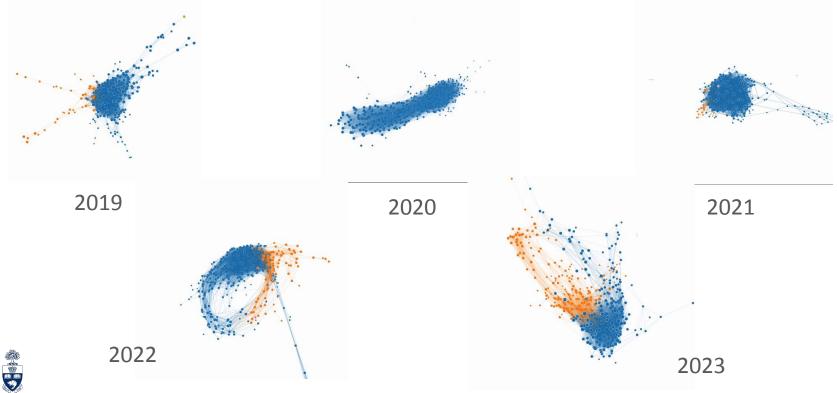
Applying Troika to Market Index Correlation Network

- Correlation Matrix
- Fisher Transform
- Thresholding
- Signed Networks

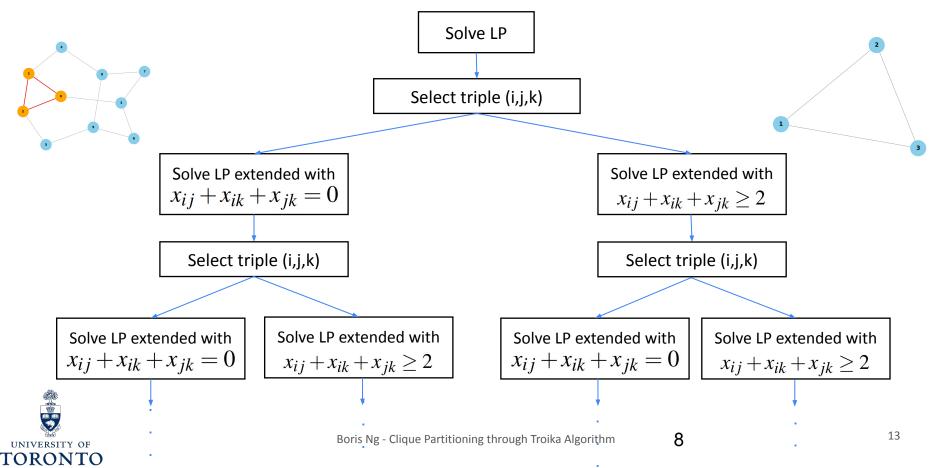




Yearly Clustering Results of S&P 500 (2019-2023)



How does Troika work? Branching on Node Triples



Pushing the computational limits

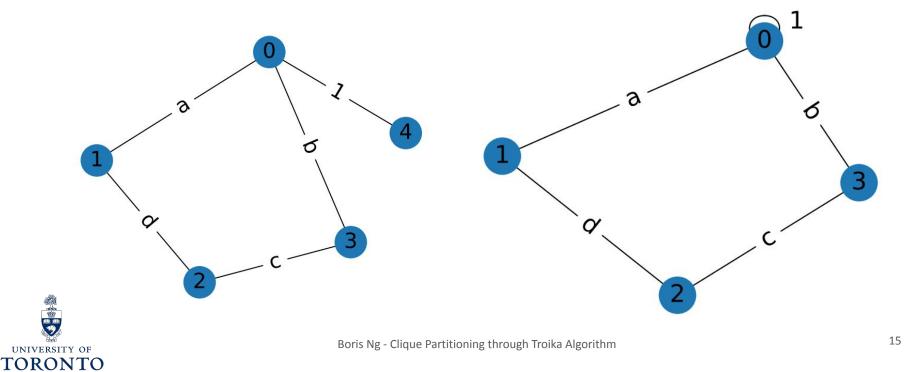
- 1. Graph Pre-processing
- g
- 2. Search Termination Criteria
- 3. Variable fixing
- 4. Implied branching
- 5. Parallel processing of the separating sets
- 6. Node Triple Selection





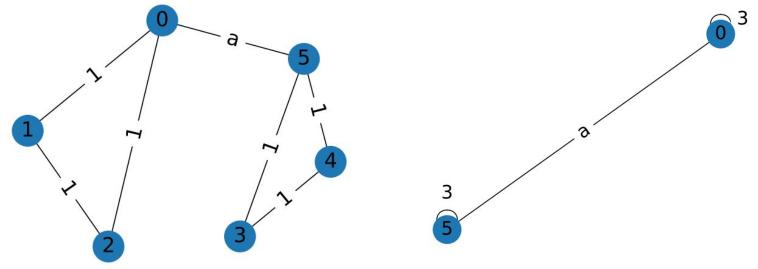
Graph Pre-processing

1. Pendant Node Reduction



Graph Pre-processing

2. Pendant Clique Reduction





Experimental Setup

- Datasets:
 - ABR Graphs: Aggregation of Binary Relations for qualitative data analysis.
 - MCF Graphs: Machine Cell Formation problems in manufacturing
 - Barabási-Albert Graphs: Scale-free networks to model real-world structures

Sorensen & Letchford (2017) CP-lib: Benchmark instances of the clique partitioning problem Barabási & Albert (1999) Emergence of scaling in random networks



Experimental Setup

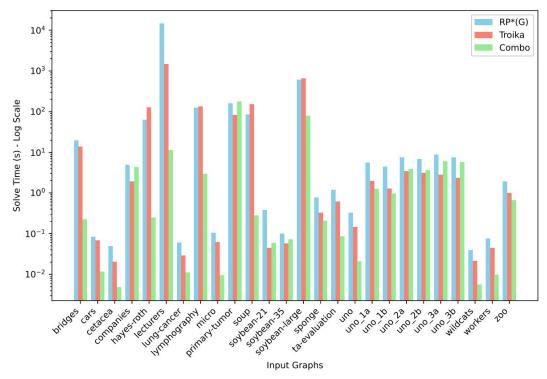
- Performance Metrics:
 - $\circ\,$ Solution Quality
 - Computational time
- Methods Compared:
 - Troika
 - $\circ~$ Gurobi IP Formulation
 - \circ Combo Algorithm

Sorensen & Letchford (2017) CP-lib: Benchmark instances of the clique partitioning problem Barabási & Albert (1999) Emergence of scaling in random networks



Solve Time

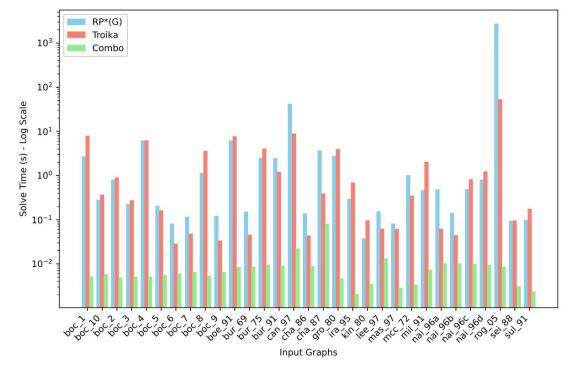






Boris Ng - Clique Partitioning through Troika Algorithm

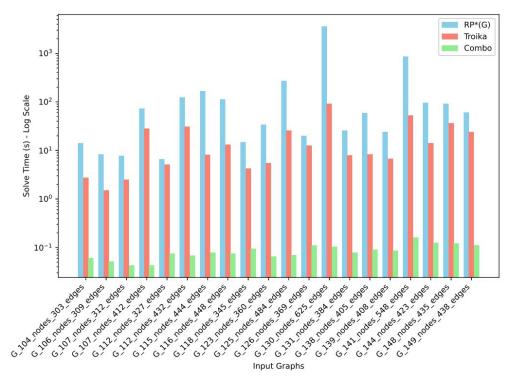
Solve Time The Machine Cell Formation (MCF) dataset





Solve Time

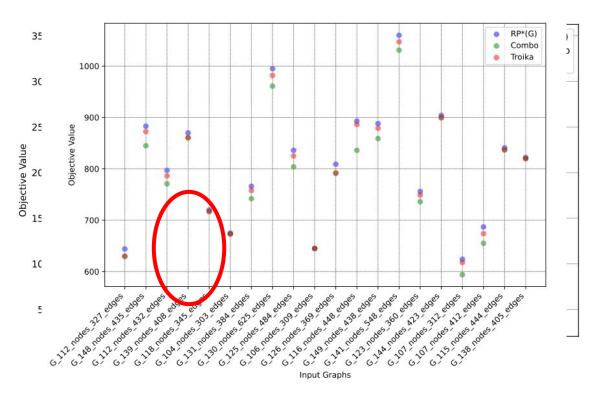
The Barabási-Albert Graphs dataset





Accuracy

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Experimental Results

• *Troika* is **5.8, 27 and 14.8 times faster** than *Gurobi IP Formulation* on average on ABR, MCF and Barabási-Albert graphs respectively

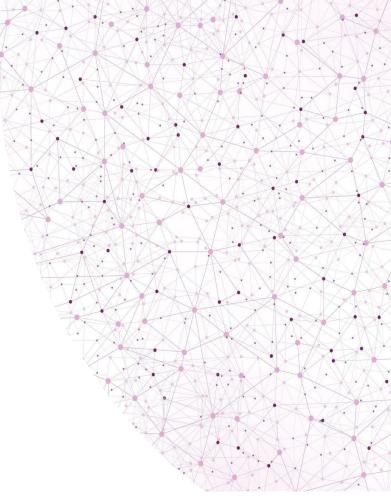
• *Troika* achieves **98.85% optimality**, compared to *Combo*'s 96.65% on both MCF and Barabási-Albert graphs.



Final Note

 Troika offers efficient and effective solution for the clique partitioning problem

 Troika demonstrates a balance in terms of solve time and solution quality





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